

Course Plan lec. 1: Introduction lec. 2: Mathematics for 3D Games lec. 3: Scene Graph lec. 4: Game 3D Physics •••• + •• lec. 5: Game Particle Systems lec. 6: Game 3D Models See also the courses lec. 7: Game **Textures** «Online Game Design» lec. 9: Game Materials «Computer Networks» lec. 8: Game **3D Animations** lec. 10: **3D Audio** for 3D Games lec. 11: **Networking** for 3D Games bridge lectures lec. 12: Rendering Techniques for 3D Games lec. 13: Artificial Intelligence for 3D Games

#### Player 2 has joined the game



- Types of multiplayer games
  - Hot-seat
    - players time-share
  - Local multiplayer (Side-to-side)
    - e.g., split screen
    - players share a terminal
  - Networked
    - each player on a terminal
    - terminals connected...
      - ...over a LAN
      - ...over the internet

Needs networking

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## Player 2 has joined the game



(see course on: Online Game Design)

- Multiplayer game types, according to gameplay
  - collaborative
  - competitive
  - versus
  - teams...
- *How much* multiplayer?
  - no: single player
  - 2 players?
  - 10 players?
  - >100?
  - > 1000?> 1000?(«massively» multiplayer online, MMO)

#### Networking in 3D Games



Objective: all players see and interact with a common 3D virtual world



how can this illusion be achieved?

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## **Networking in Games**



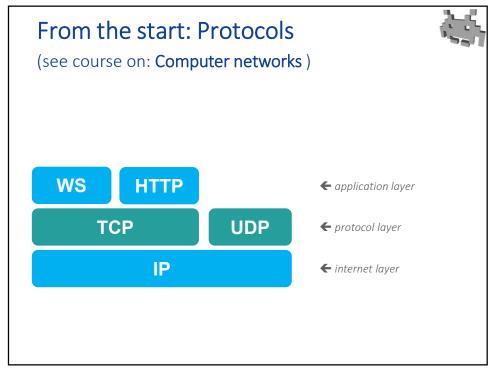
- One task of a Game Engine
- Scenarios can be very different:
  - number of players? (2, 10, 100, 100.000?)
  - game pace? (real time action ≠ chess match)
     Latency tolerance?
  - joining ongoing games : allowed?
  - cheating : must it be prevented?
  - security: is it an issue? (e.g. DoS attacks)
  - channel: LAN only? internet too? Bandwidth needed?

# Dev choices for a networked-game



- What to communicate?
  - e.g.: complete statuses, status changes, inputs...
- How often?
  - at which rate
- Over which protocol?
  - TCP, UDP, WS ...
- Over which network architecture?
  - Client/Sever, Peer-To-Peer
- How to deal with networking problems
  - latency ("lag") <== one main issue
  - limited bandwidth
  - connection loss
  - loss of packets

/



#### **Protocol layer**



#### **TCP** sockets

- Connection based
- Guaranteed reliable
- Guaranteed ordered
- Automatic breaking of data into packets
- Flow control
- Easy to use, feels like read and write data to a file

#### **UDP** sockets

- What's a connection?
- No reliability
- No ordering
- Break your data yourself
- No flow control
- Hard.
   Must detect and deal with problems yourself.

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#### **UDP vs TCP**



Nagle's

algorithm

caching?

no, thank you

- The problem with TCP: too many strong guarantees
  - they cost a lot in terms of latency ("lag")
  - not designed for time-critical applications
  - if it must be used, at least enable the option TCP NODELAY
- The problem with UDP: not enough guarantees
  - no concept of connection: no timeouts, no handshake, a port receives from anyone
  - List of guarantees:
     1. "a packet arrives either whole, or not at all".
     End of list.
  - packets can arrive...
     ...out of order :-O, ...not at all :-O, ...in multiple copies :-O

#### **UDP vs TCP**



- Problem with TCP
  - too many costly guarantees
- Problem with UDP
  - not enough guarantees
- The hard way:
  - use UDP, but manually re-implement a few guarantees



- best, for the most challenging scenario
  - fast paced games, not on LAN

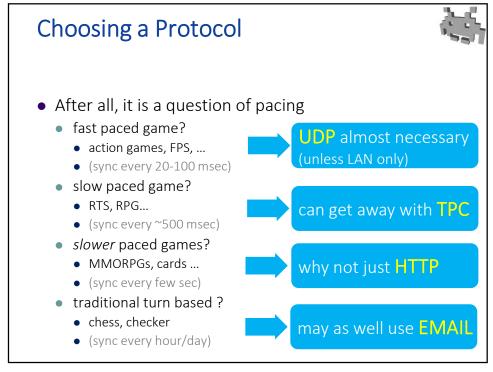
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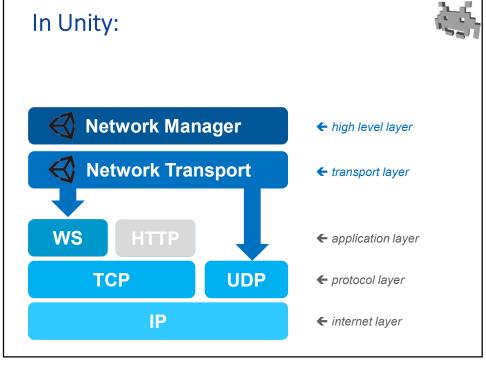
# Virtual connections over UDP: how-to (notes)



- add connection ID to packets
  - to filter out unrelated ones
- time out on prolonged silence (~ few secs)
  - declare "connection" dead
- add serial number to packets
  - to detect when one went missing / is out of order / is duplicate
  - (warning: int numbers do loop solutions?)
- give ack back for received packets
  - optimize for lucky (& common) cases!
    - N (say 100) received msg == 1 ack (with bitmask)
  - resend? only a few times, then give up (data expired) <
- congestion avoidance: measure delivery time
  - tune send-rate (packets-per-sec) accordingly
- obviously: NON blocking receives!

what **TPC** doesn't understand





#### In Unity:



- Low level: Transport Layer
  - Implements guarantees over UDP (virtual connections)
  - As easy to use as TCP, but designed for games
    - see how-to notes list above
  - Can work over WS instead UDP (abstracts the differences)
    - WS is used for web games (the one with graphics on WebGL)
- Hi level: Network Manager
  - presets network connectivity
  - by default: "client hosted" games
    - the server is one of the players
  - controls shared state of the game
  - deals with clients
  - sends remote commands

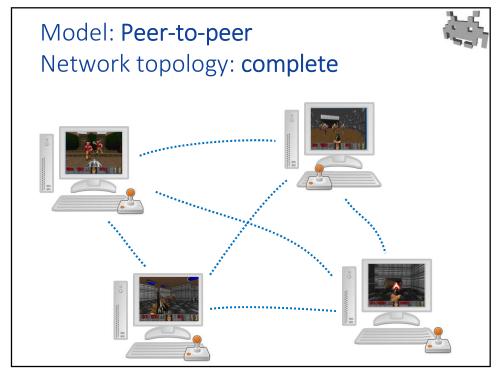
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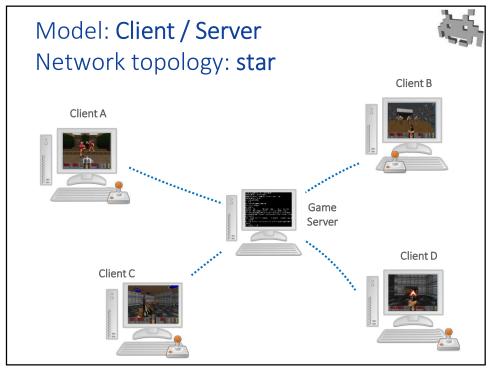
## Networking paradigms for games: we will see...

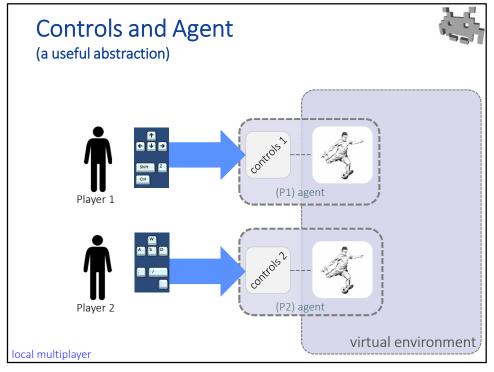


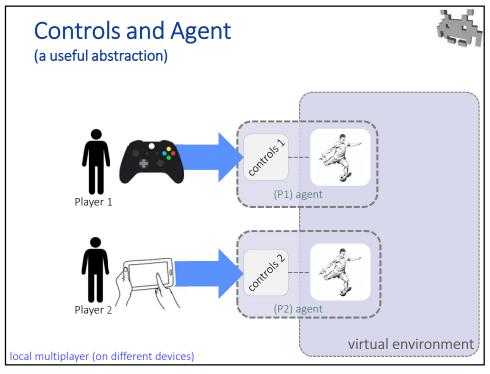
- Deterministic Lockstep on P2P
- Deterministic Lockstep
- Game-Status Snapshots
- Distributed Physics (just notes)
- Game-Status Snapshots with Client-Side predictions
- Cloud gaming

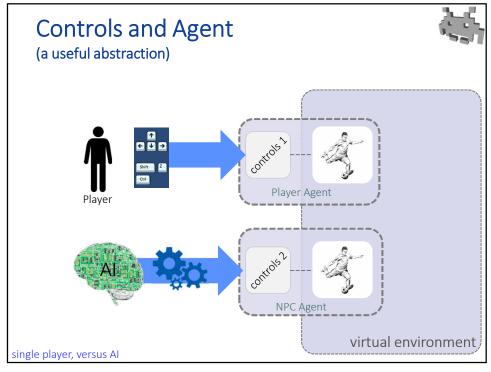
on client-server

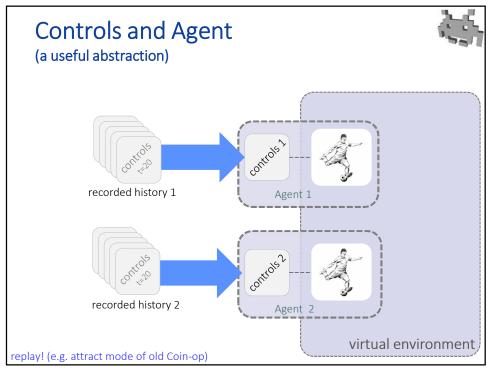


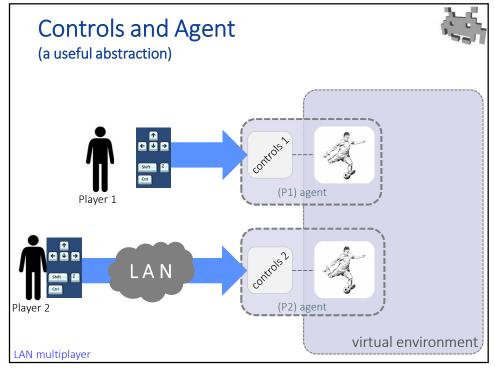


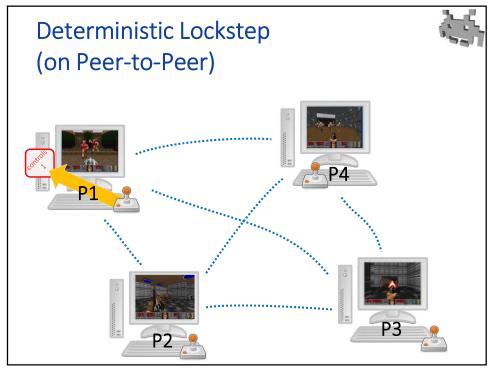


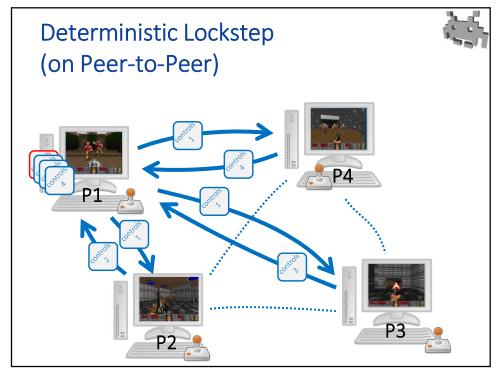


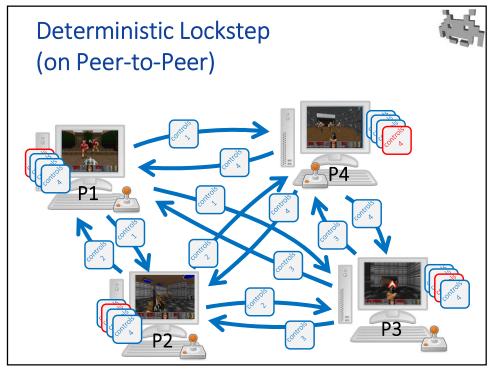


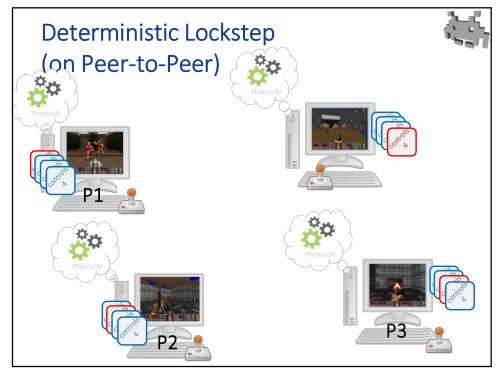












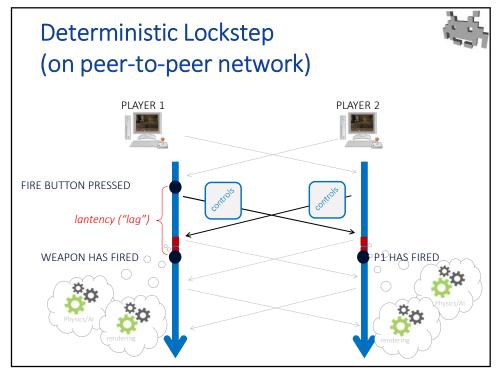
# Deterministic Lockstep (on Peer-to-Peer)





- Game evolution = sequence of "turns"
  - e.g. physics steps (fixed dt!)
- Each node sends its current controls (inputs)
  - to everybody else
- After all controls are received, each node computes its own evolution
  - deterministically: same input → same result





## Deterministic Lockstep: the good



- elegant and simple! ©
- minimal bandwidth needed
  - only sent data = controls
    - compact! (e.g., a bitmask)
  - does not depend on complexity of virtual environment
- cheating: inherently limited
  - but a few ways to cheat are still possible, e.g.:
    - aim-bots (unlawful assist from AI)
    - x-rays (unlawful reveal of info to player)
- mixes well with:
  - non-cheating AI, replays, player performance recording...
- can use simple TCP connections
  - because we need 0% packet loss anyway (but...)

#### **Deterministic Lockstep:** can as well use TPC instead of UDP?



- why yes:
  - TPC is simple to use
    - It takes care of everything
  - works well, when no packet loss
    - on loss, we need resend it anyway: let TPC do that
  - makes little sense to use UDP and then... try to re-implement all TPC over it
  - at the beginning of dev, **UDP** is a (premature) optimization
- why not:
  - to degrade better with packet loss
  - e.g.: use redundancy instead of resend-on-failure
    - controls are small: send 100+ controls in every packet
    - keep resending until ack received

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## **Deterministic Lockstep**



- Common, e.g., in:
  - RTS
    - controls = orders
      - can be fairly complex
      - but game status = much more complex
  - first generation FPS
    - controls = [gaze dir + key status]















Ensemble Studios et al. 1998..2015



Doom ID-soft, 1998

## **Deterministic Lockstep** on peer-to-peer: the limitations



- responsiveness:
  - input-to-response delay of 1 x delivery time (even locally!)
  - (you cannot act immediately even on your own local input)
- does not scale with number of players
  - quadratic number of packets
  - 2P ok, 100P not ok
- input rate = packet delivery rate
- delivery rate = as fast as the slowest connection allows
- if connection problems (anywhere): everybody freezes!
- joining ongoing games: difficult to implement
  - needs sends full game state to new player
- assumes full agreement on initial conditions
  - this is not problematic
- assumes complete determinism!
  - this can be problematic (see next slide)

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#### Determinism: what can break it



The entire game system

"determinism" in mind ...

...and still, it difficult to get

must be designed

from the start with

(and debug)

- Pseudo-Random? → not as problem
  - fully deterministic (just agree on the seed)

**Physics**: many preclusions and traps ▲ variable time step? bad

▲ time budgeting? bad

▲ hidden threats:

order of processing of particles/constraints

A anything that depends on clock?

→ poison to determinism

**A** GPU computations? very dangerous

slightly different outcome on each GPU model

**floating point** operations?

hidden dangers, different hardwired implementations

best to assume very little (fixed point is 100% safe)

NOTE: 99.999% correct == not correct

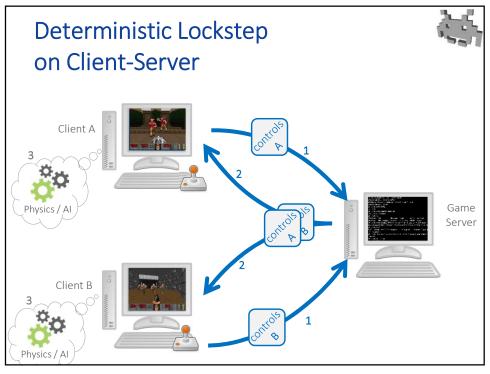
virtual world is faithful to reality enough to be *chaotic* → butterfly effect: the tiniest local difference == expect completely different outcomes soon

# A limitation of any peer-to-peer paradigm



- They do not scale with number of players
  - quadratic number of packets necessary
  - 2Players ok, 100Players not ok
- Every client needs to send/receive from all other clients
  - Topology of the network: fully connected
  - E.g. if implemented with TCP: tons of connections
  - Not always practical (firewalls, etc)
- Let's switch to Client / Server paradigms (topology of the network: star)

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## Deterministic lockstep on Client-Server



- Server sits on the central node
- Protocol:
  - Each client sends his controls to server
  - Server collects all controls and sends them back to clients
- Advantage:
  - scalability: number of packets is linear (not quadratic)
- Cost:
  - responsiveness:latency = 2 × delivery time \_\_:-0

hurts gameplay!

- Bonus: the server can now be made authoritative
  - Many new options available. For example...

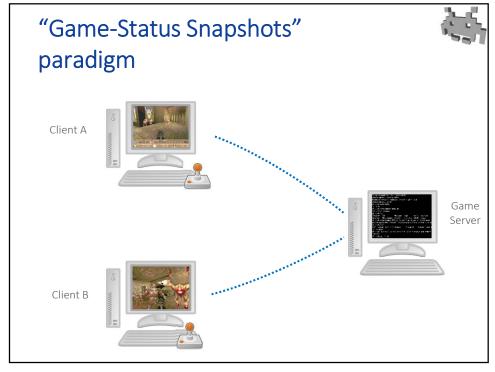
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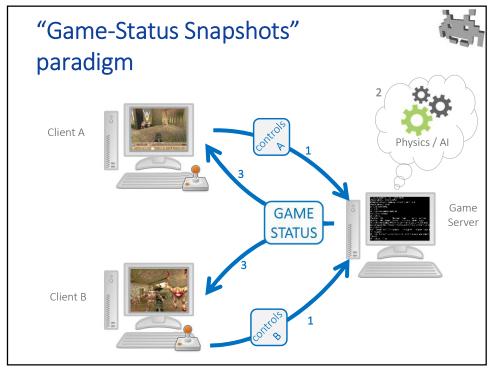
# "Server is the man" \* (authoritative server ). Concept:

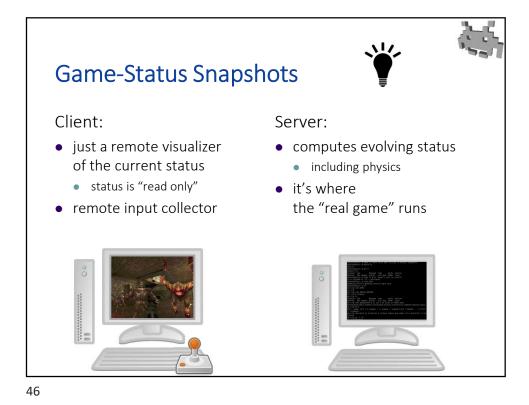


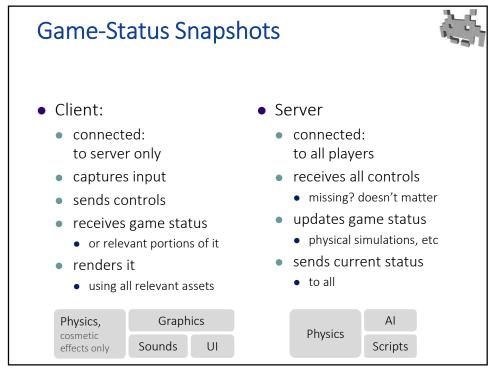
- The server has the last word
- For example:
  - Packet loss from player 3?
     Server makes up controls for player 3 (instead of waiting for them)
  - Note: server defines what player 3 did, not player 3 itself!
  - i.e., clients take server's word even for their own actions
  - Packet loss affects one player only

\* Tim Sweeney (Unreal / Epic Games)









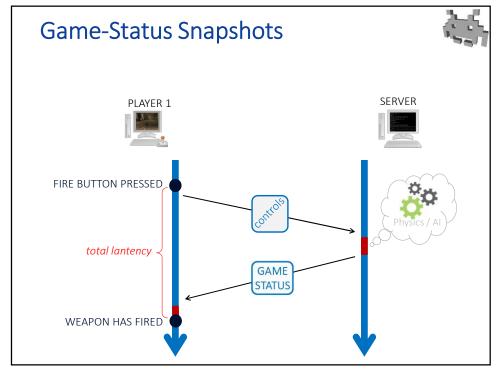
## Game-Status Snapshots (compared to deterministic lockstep)

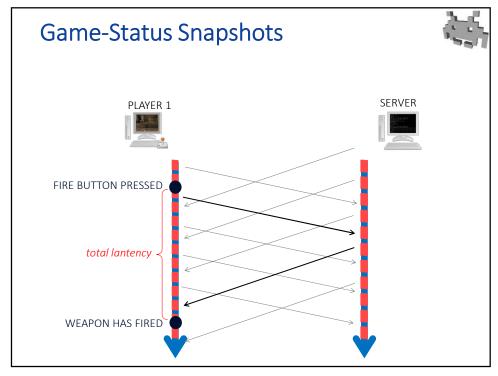


- the gains:
  - determinism: not a required assumption anymore
  - joining ongoing games: becomes trivial
  - packet loss: bearable (hurts that player only)
    - to profit: use UDP
  - slower connection: bearable (affects that player *only*)
- the losses:
  - packet size: a lot bigger! optimizations, to counter this:
    - compress world status
    - send only the portions of the status which changed or which interest a player
  - perceived latency: from input to effect = delivery time :-( from input to visual = 2 x delivery time :-0

Loss of responsiveness hurts gameplay!

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## Game-Status Snapshots: with Interpolation: the idea



- World "Snapshot" contains:
  - The scenegraph (local transforms per object)
  - Anything else needed e.g. for gameplay
- Problems:
  - snapshot size can be large! (even with optimizations)
  - few FPS (in the physical simulation)
  - ==> "jerky" animations
- Solution 1: client-side interpolation
  - client keeps last two snapshots in memory
    - last received one + the previous one
  - interpolates between them,
    - client lags behind server by even more!
  - gain: smoothness (high FPS with low packet rate)
  - loss: responsiveness (increased latency) oh noes!

## Game-Status Snapshots: with Extrapolation: the idea



- World "Snapshot" contains:
  - data needed for 3D rendering: (position-orientation of objects, plus anything else needed)
- Problems
  - large snapshot size! (even with optimizations)
  - few FPS (in the physical simulation)
  - ==> "jerky" animations
- Solution 2: client-side extrapolation
  - clients keeps last two snapshots in memory
    - last received one + the previous one
  - extrapolates between them, i.e., shows the expected "future"
    - i.e. it shows an attempted prediction to the next snapshot
    - NOTE: this prediction is often wrong: glitches.
  - gain: responsiveness
  - loss: accuracy lots of glitches. :-(

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# Partial Client-side Game Evolution (aka distributed physics): the idea



to server,

or, in a P2P network,

to each other peers

- Each client:
  - in charge for game evolution
    - including physics
  - communicates to others a reduced game-status snapshot
    - describes only status of own player (e.g. it's transforms, transforms for its flying bullets, etc)
  - receives other partial snapshots
  - merges everything up
    - (updates statuses of other players)
- Simple, zeroed latency
  - immediately responsive to local player controls
  - remote agents updated according to "what their client says"
- Problem: can still need determinism
  - (who keeps NPCs / environment in sync?)
- Problem: authoritative clients: prone to cheating!!!

## Client-Side Prediction: the idea





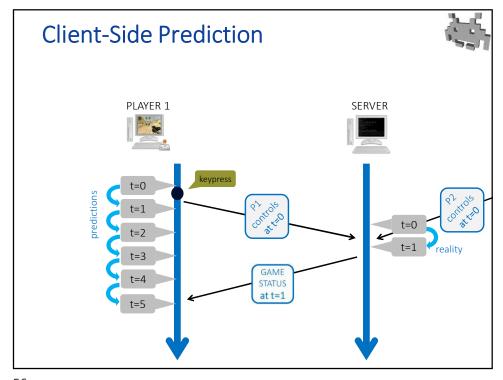
- Client:
  - get Commands from local inputs
  - sends Commands to Server
  - computes game evolution (the prediction)
    - maybe "guessing" other players commands (which it ignores)
    - zero latency!
- Server:
  - receives Commands (from all clients)
  - computes game evolution (the "reality")
    - server is authoritative
    - prevents many forms of cheating
  - sends Snapshot back (to all clients)
- Client:
  - receives Snapshot (the "real" game status)
  - corrects its prediction, only if needed

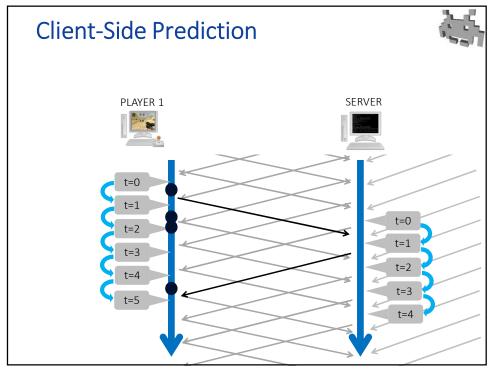
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## Client-Side Prediction with corrections from the server



- The server-side "real" simulation lives h msecs in the past of the client-side "predicted" one
  - h = deliver time
  - remember: virtual time != real world time
- When server correction arrives to client, it refers to 2h msecs ago (for the client)
- Q: how to correct... the past?





## Client-Side Prediction: correction from the server



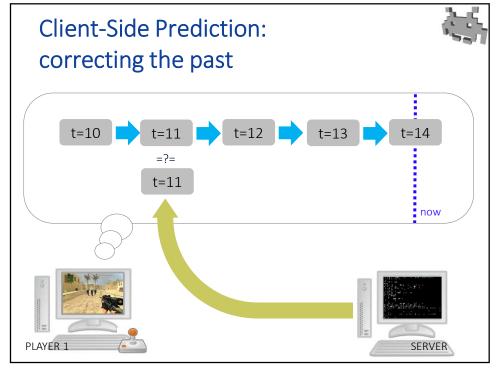
- Q: How to correct... the past?
- A
  - keep last N statuses in memory
    - including own controls
  - as the "real" status (the correction) of the past arrives from server...
  - ...compare it with stored past status (at corresponding time):
    - does it match? within some tolerance
    - does it mismatch? discard frame and following ones, rerun simulation to present (reusing stored controls)

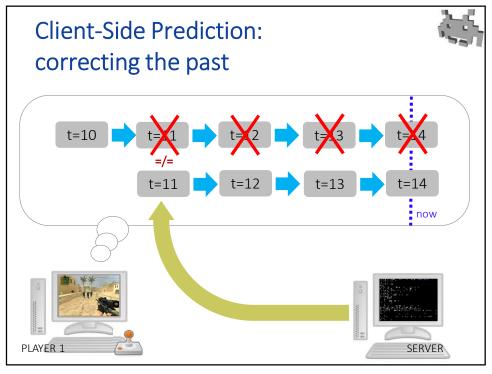
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#### Re-running physical simulation



- We just need to catch up with the present
- Physics and AI only
  - no graphics, no sound rendering, no cosmetic physics, particle system...
- At full speed: can use larger dt if necessary
  - This only compromises accuracy a bit -- tolerable
- Must reuse same controls of own player and other's
  - They must all be cached also
- Note: player is never shown these intermediate steps; only the final result
- The price to be paid: glitches when switching from current present to a different (corrected) present





# Client-Side Prediction: what causes mispredictions?



commonest

- Lack of determinism
  - e.g., physics was approximated *soft* real-time!
  - see above for more possible causes of this
- Didn't account that own controls were not received by server (in time)
  - server: "actually, back then, you didn't jump"
  - authoritative server server defines the truth (even when the client is in a better position to know)
- Didn't account for other players' controls
  - server: "actually, back then, that other player jumped"
- Note: none of the above breaks gameplay too much
  - it just causes minor / temporary glitches (usually)

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# Client-Side Prediction: optimizations 1/2



- reduce snapshots size
  - (==> to increase packet frequency)
  - partial snapshots: refresh more often the parts which are most likely to be predicted wrong / or which changed
  - drastic space reductions!
  - but make sure that every part is eventually refreshed
- reduce correction computation

(==> making corrections quicker)

- partial physic steps: update only the parts affected by the error
- use bigger dt (fewer steps to get to present)

## Client-Side Prediction: optimizations 2/2



- tentatively predict also unknown data
   (==> to reduce correction frequency)
  - e.g., also predict other player's controls
  - easiest prediction: players do what they did last frame
- trigger correction only when status differ enough
   (==> to reduce correction frequency)
  - e.g., when any spatial position difference > epsilon
  - tolerate small discrepancies
  - (warning: discrepancies tend to explode exponentially with virtual time because Chaos)

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## Client-Side Prediction: notes



- A snapshot = includes physical data
  - not just for the 3D rendering, also to update physics
  - can be made smaller, with optimizations
- ② No latency: immediately react to local input
  - client proceeds right away with next frame
  - when prediction is correct: seamless illusion
  - otherwise: (minor?) glitches
- © Determinism: not assumed
- © Cheating: not easy (server is authoritative)

#### Summary: rules of thumb



- Comparing network layouts
  - peer-to-peer :
    - © reduced latency
    - • quadratic number of packages
       (with number of players)
  - client-server :

    - ② linear number of packages (with number of players)
    - REQUIRED, for any solution with a authoritative server
    - REQUIRED, for num players >> 4-6

preventing cheating not important

• game status not overly complex

Client-side prediction + server correction, if

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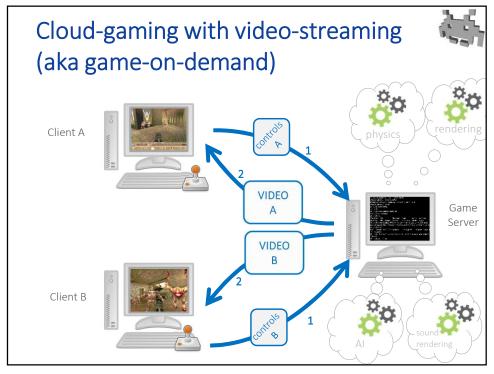
# Comparing network paradigms Deterministic Lockstep, if determinism can be assumed few players (up to 4-5) fast + reliable connection (e.g., LAN) Game-status Snapshots, if game status not overly complex a little latency can be tolerated Client-side evolution, if

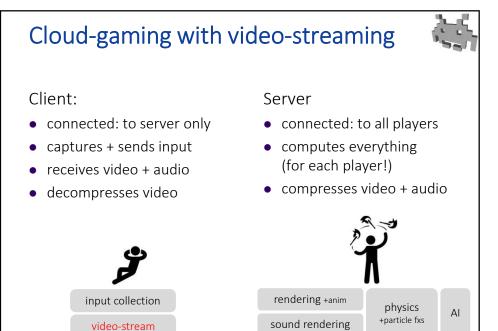
## Summary: classes of solutions



- Who computes game evolution? (including physics)
  - deterministic-lockstep : clients
    - there may be no server at all: peer-to-peer
    - independent computations, same result
  - game-status snapshots : server
    - clients are just visualizers
    - maybe with interpolation / extrapolation
  - (distributed physics : both clients and server)
    - clients in charge for own agent(s)
    - server in charge for env. / NPCs
  - client-side predictions : both clients and server
    - clients "predict" (just for local visualization purposes)
    - server "corrects" (it has the last word!)

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# Cloud-gaming with video-streaming (aka game-on-demand)



- Compared to game snapshots technique...
- Server now does everything
  - 3D Physics + AI (same as with Game-Status snapshots)
  - 3D Rendering (aka "remote" rendering) (including animations, particle effects...)
  - 3D Sound rendering

decompression

- Compresses and send 2D-video/audio
- ×N times! (once for each player)

video-stream compression

- Client does almost nothing
  - Collects and sends inputs (controls)
  - Receives and decompresses 2D-video/audio

## Cloud-gaming with video-streaming (aka game-on-demand)



- Advantages: client is thin
  - client does almost nothing
  - client needs nothing (no asset, no storage, no DRM)
  - needed capabilities are limited (pads, cellphones ok)
- Challenges:
  - Demanding in terms of bandwidth (high-res video + audio)
  - Demanding in terms of server workload
  - Latency! Impossible to reduce or hide (by prediction), plus compression by server, luckily, video-on-demand plus decompression by client technologies can be leveraged
  - Video resolution, FPS: are very costly

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## Cloud-gaming (aka gaming-on-demand)





• A heavily invested-on, « fast-growing approach to 3D game networking



Latency = maybe 80-120 ms



• Is this acceptable?



Bandwidth = maybe 5-50 mbits/s



 Will it become an established model for online 3D games?